

Type Definitions (from Gamut vol 2, pp. 81-82)

1. If α is a variable or a constant of type a in L , then α is an expression of type a in L .
2. If α is an expression of type $\langle a, b \rangle$ in L and β is an expression of type a in L , then $(\alpha(\beta))$ is an expression of type b in L .
3. If φ and Ψ are expressions of type t in L (i.e. formulas in L), then so are $\neg \varphi$, $(\varphi \wedge \Psi)$, $(\varphi \vee \Psi)$, $(\varphi \rightarrow \Psi)$, and $(\varphi \leftrightarrow \Psi)$.
4. If φ is an expression of type t in L and v is a variable (of arbitrary type a) then $\forall v\varphi$ and $\exists v\varphi$ are expressions of type t in L .
5. If α and β are expressions in L which belong to the same (arbitrary) type, then $(\alpha = \beta)$ is an expression of type t in L .
6. Every expression in L is to be constructed by means of (1)-(5) in a finite number of steps.

Restatement and Extension to Intensional Types (pp. 119-120)

1. If $\alpha \in \text{VAR}_a$ or $\alpha \in \text{CON}_a^L$, then $\alpha \in \text{WE}_a^L$.
2. If $\alpha \in \text{WE}_{\langle a, b \rangle}^L$ and $\beta \in \text{WE}_a^L$, then $(\alpha(\beta)) \in \text{WE}_b^L$.
3. If $\varphi, \Psi \in \text{WE}_t^L$, then $\neg \varphi$, $(\varphi \wedge \Psi)$, $(\varphi \vee \Psi)$, $(\varphi \rightarrow \Psi)$, and $(\varphi \leftrightarrow \Psi) \in \text{WE}_t^L$.
4. If $\varphi \in \text{WE}_t^L$ and $v \in \text{VAR}_a$, then $\forall v\varphi$ and $\exists v\varphi \in \text{WE}_t^L$.
5. If $\alpha \in \text{WE}_a^L$ and $v \in \text{VAR}_b$, then $\lambda v\alpha \in \text{Var}_{\langle b, a \rangle}$.
6. If $\varphi \in \text{WE}_t^L$, then $\Box\varphi, \Diamond\varphi \in \text{WE}_t^L$.
7. If $\alpha \in \text{WE}_a^L$, then $\hat{\alpha} \in \text{WE}_{\langle s, a \rangle}^L$.
8. If $\alpha \in \text{WE}_{\langle s, a \rangle}^L$, then $\forall\alpha \in \text{WE}_a^L$.
9. Every element of WE_a^L for any a is constructed in a finite number of steps using (1)-(9).